

All-to-All Communication

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Principles of Distributed Computing

Overview



- I. Introduction
- II. Previous Results
- III. Fast MST Algorithm
- IV. Analysis
- V. Summary
- VI. Extensive Example



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Overview



- I. Introduction
 - Definitions & System Model
 - A Simple Algorithm
- II. Previous Results
- III. Fast MST Algorithm
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I. Introduction



Definitions

- A **tree** is a **connected** graph **without cycles**.
- A **subgraph** that **spans all** vertices of a graph is called a **spanning subgraph**.
- Among all the **spanning trees** of a **weighted** and **connected graph**, the one (possibly more) with the least total weight is called a **minimum spanning tree (MST)**.



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I. Introduction



Definitions

- In a **MST algorithm**, $|V| - 1$ edges have to be chosen in total. In each **phase** of the algorithm, probably only a fraction of those edges are chosen.
- Nodes that are directly or indirectly connected using chosen edges only belong to the same **cluster**.
- The **minimum weight outgoing edge (MWOE)** is the edge with the lowest weight among all incident edges leading to other **clusters**.



I. Introduction



Usage of MST

Minimize the cost associated with global operations such as **broadcasts!**

→ Minimize the message complexity:

Avoid traffic explosion by using a **spanning tree** (no cycles!)

→ Minimize the time complexity:

If the edge weights represent the delay on the link, then MST minimizes the **execution time**.

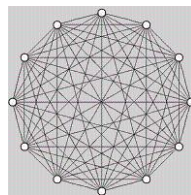


I. Introduction



System Model

- The system is represented by a **complete weighted undirected graph**
- $G=(V,E,w)$ where $w(e)$ denotes the weight of edge $e \in E$ and $|V|=n$. All **edge weights** are different (w.l.o.g.).
- Each node has a **distinct ID** of $O(\log n)$ bits.
 - Each node knows all the edges it is incident to and their weights.
 - Each node knows about all the other nodes.
 - The **synchronous communication model** is used.



I. Introduction



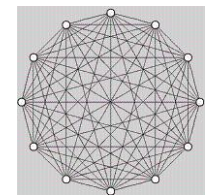
Synchronous Communication Model

Communication advances in **global rounds**.

In each **round**, processes **send** messages, **receive** messages and do some **local computation**.

The **time complexity** is the number of rounds until the computation terminates in the worst case.

The **message complexity** is the number of messages exchanged in the worst case.





Getting a Feeling for the Problem...

How hard is it to compute the MST in a distributed system (assuming a fully connected graph)?

All nodes know the weights of all incident edges. If all nodes send this information to all other nodes, then all nodes suddenly have the entire picture!

→ A simple algorithm that requires only one round!

However, that is not really interesting...

Therefore, the message size is limited to $O(\log n)$ bits!

Note that the previous algorithm requires messages of size $O(n \log n)$!



Getting a Feeling for the Problem...

Since each **node ID** (and edge weight) requires $O(\log n)$ bits, this implies that only a constant number of node IDs (and edge weight) can be packed into a single message!

We demand that all nodes know the **MST** at the end of the computation!

How can the MST be constructed now?

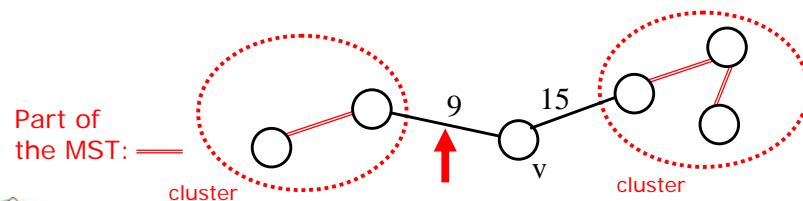
→ Let's look at a simple algorithm first...



A Simple Algorithm

All MST algorithms (local or distributed) are based on the following lemma:

Lemma 1: It is always safe to add an edge to the spanning tree, if this edge is the **MWOE** of a node v or a cluster F .



A Simple Algorithm

Phase k: Code for node v in cluster F

Input: Set of chosen edges that build node clusters

1. Compute the MWOE
2. Send the MWOE to all nodes in the same cluster
3. Receive messages from the other nodes
4. If own MWOE is the lightest, then broadcast it to all other nodes and add this edge (→ All nodes have to know all clusters after each round)
5. Receive other broadcast messages and add those edges as well



I. Introduction of the Problem

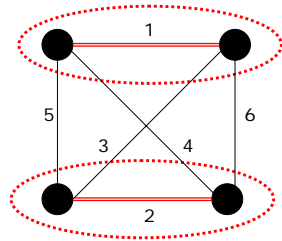
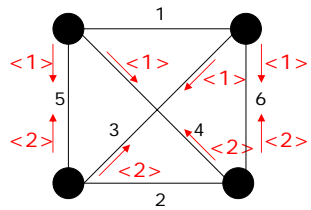


A Simple Algorithm

Example

$$\langle w(\{v,u\}) \rangle = \langle v,u,w(\{v,u\}) \rangle$$

Round 1:



Broadcast the lightest edge to the other nodes

Add edges and update clusters



I. Introduction

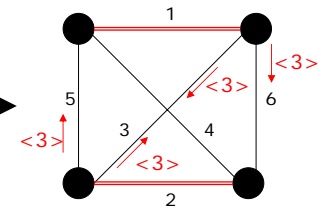
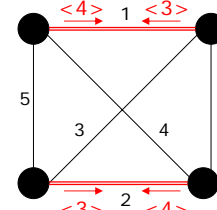


A Simple Algorithm

Example

$$\langle w(\{v,u\}) \rangle = \langle v,u,w(\{v,u\}) \rangle$$

Round 2:



Send MWOE to all nodes in the same cluster

Broadcast the lightest edge to the other nodes



I. Introduction

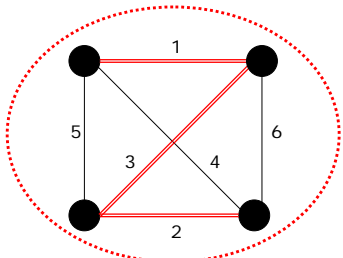


A Simple Algorithm

Example

$$\langle w(\{v,u\}) \rangle = \langle v,u,w(\{v,u\}) \rangle$$

Round 2:



Add edges and update clusters

The algorithm is obviously correct. Since the minimum cluster size doubles in each round, the algorithm computes the MST in $O(\log n)$ rounds!

Can it be improved???

Lower bound???



Overview



- I. Introduction
- II. Previous Results
 - Lower and Upper Bounds
 - Open Questions
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II. Related Work

Previous Results

D denotes the **diameter** of the graph, i.e., the maximum **distance** between any two nodes of the graph.

	Known Algorithms	Known Lower Bounds
$D = 1$	$O(\log n)$???
$D = 2$	$O(\log n)$???
$D \geq 3$	$O(D + \sqrt{n} \log^* n)$	$\Omega(n^{1/4}/\sqrt{\log n})$



II. Related Work

Previous Results

Even our simple algorithm has this complexity!!

Interesting „jump“!

	Known Algorithms	Known Lower Bounds
$D = 1$	<u>$O(\log n)$</u>	???
$D = 2$	$O(\log n)$???
$D \geq 3$	$O(D + \sqrt{n} \log^* n)$	$\Omega(n^{1/4}/\sqrt{\log n})$



II. Related Work

Previous Results

We will now derive an algorithm with time complexity $O(\log \log n)$!

	Known Algorithms	Known Lower Bounds
$D = 1$	<u>$O(\log \log n)$</u>	???
$D = 2$	$O(\log n)$???
$D \geq 3$	$O(D + \sqrt{n} \log^* n)$	$\Omega(n^{1/4}/\sqrt{\log n})$



Overview

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 - General Idea & Problems
 - The Algorithm Step by Step
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III. Fast MST Algorithm



General Idea

In order to reduce the number of rounds, obviously clusters have to grow faster!

In our simple algorithm, we used the **MWOE** of each cluster to merge clusters.

With this approach, the **minimum cluster size doubled** in each phase.

It would certainly be faster, if the **k lightest outgoing edges** of each cluster were used, where **k** is the number of nodes in the cluster!

→ This is exactly what our new algorithm will do!



III. Fast MST Algorithm



General Idea

In order to reduce the number of rounds, obviously clusters have to grow faster!

Let β_k denote the **minimum cluster size in phase k**, then it holds for our simple algorithms that

$$\beta_{k+1} \geq 2 \cdot \beta_k$$

and

$$\beta_0 := 1$$

thus

$$2^k \leq \beta_k \leq n \rightarrow k \leq \log_2 n$$



III. Fast MST Algorithm



General Idea

Our goal is to improve the following inequality:

$$\beta_{k+1} \geq 2 \cdot \beta_k$$

We will derive an algorithm for which it holds that:

$$\beta_{k+1} \geq \beta_k (\beta_k + 1)$$

Thus the cluster sizes grow quadratically as opposed to merely double in each phase! In order to achieve such a rate, information has to be spread faster!

We will use a simple trick for that...

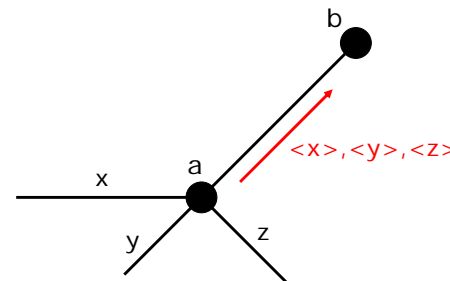


III. Fast MST Algorithm



General Idea

Unfortunately, we cannot send a lot of information over a single link...



III. Fast MST Algorithm



General Idea

However, we can send a lot of information from different nodes to a particular node v_0 !

A node can simply send parts of the information that it wants to transmit to a specific node to other nodes. These nodes can send all parts to the specific node in one step!!

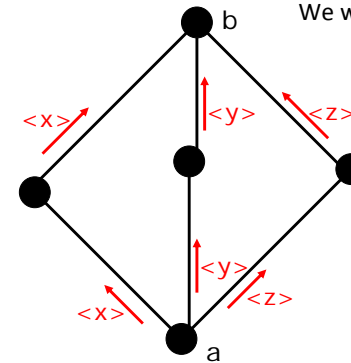
This can be used to share workload!!!



III. Fast MST Algorithm



General Idea



We will use this trick twice in our algorithm!



III. Fast MST Algorithm



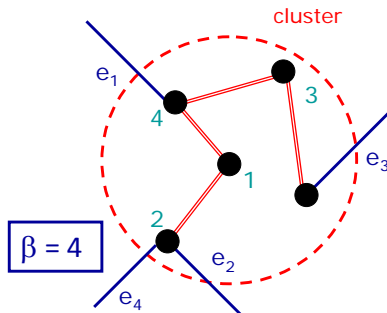
General Idea

Our new algorithm will execute the following steps in each phase.

Let β be the minimal cluster size.

1. Each cluster computes the β lightest edges e_1, \dots, e_β to other distinct clusters

2. Assign at most one of those lightest edges to the members of the cluster!



III. Fast MST Algorithm

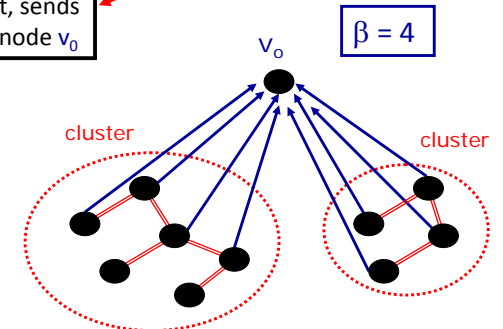


General Idea

3. Each node with an edge $\langle v, u, w(\{v, u\}) \rangle$ assigned to it, sends $\langle v, u, w(\{v, u\}) \rangle$ to a specific node v_0

Step 2 and 3 together is exactly our trick!

4. Node v_0 computes the lightest edges that can be safely added to the spanning tree



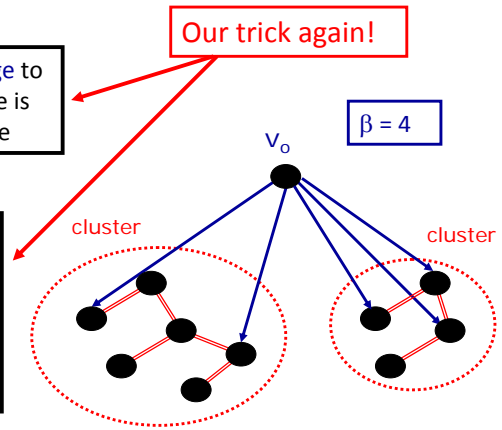
III. Fast MST Algorithm



General Idea

5. Node v_0 sends a message to a node, if its assigned edge is added to the spanning tree

6. Each node that received a message broadcasts it to all other nodes (\rightarrow All nodes have to now about all added edges!)



III. Fast MST Algorithm



General Idea

This way, more edges can be added in one phase!
However, it is not clear yet how fast it really is...

Furthermore, we do not know yet how these steps work in detail!!!
 \rightarrow There are a few obvious problems...



III. Fast MST Algorithm



Problems

First problem:

How can the β lightest outgoing edges of a specific cluster be computed?

\rightarrow This is actually not so difficult. The procedure [Cheap_Out](#) in the actual algorithm copes with this problem. We will discuss it in the following section.



III. Fast MST Algorithm



Problems

Second problem:

How can the designated node v_0 know which edges can be added without creating a cycle in the constructed graph?

Let's illustrate the problem with an example graph!



III. Fast MST Algorithm



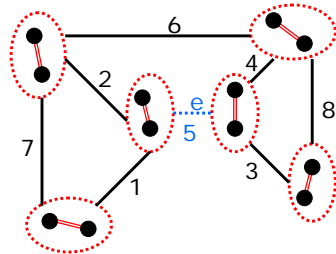
Problems

In our example:

$|V| = n = 12$

$\beta = 2$ (minimum cluster size)

This is the picture of the designated node v_0 after receiving the $\beta = 2$ lightest outgoing edges of each cluster



v_0 does not know about the edge e ! It is the 3rd lightest edge of both adjacent clusters!

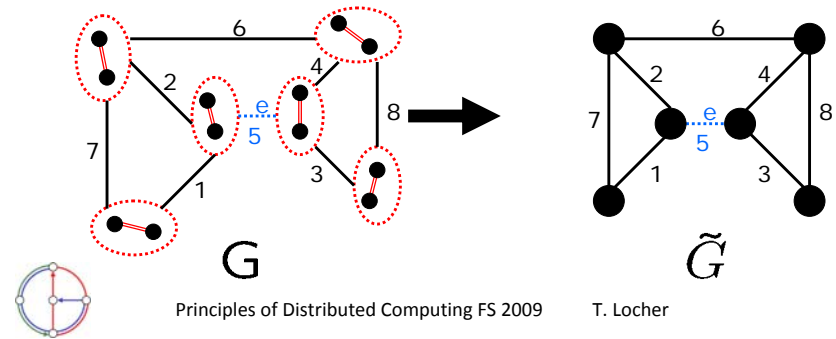


III. Fast MST Algorithm



Problems

v_0 can construct a logical graph. Its nodes are the clusters and its edges are the $\beta = 2$ lightest outgoing edges.

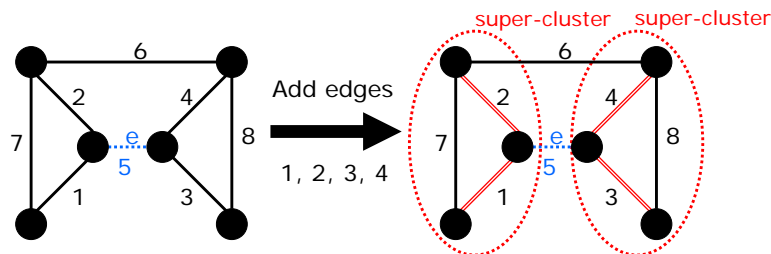


III. Fast MST Algorithm



Problems

Based on the knowledge of the $\beta = 2$ lightest outgoing edges, v_0 can locally merge nodes of the logical graph into clusters. The 4 edges with weights 1, 2, 3 and 4 can be chosen safely, since always the MWOE is used.



III. Fast MST Algorithm



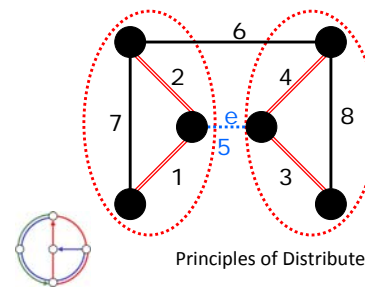
Problems

If the edge with weight 6 is used to finish the construction of the spanning tree, then the resulting tree is not the MST!!!



The problem is that in both (super-) clusters, at least one of the clusters has already used up all of its β outgoing edges. The $(\beta+1)$ th outgoing edge might be lighter than other edges!!!

So, when is it safe to add an edge???



III. Fast MST Algorithm



The Algorithm Step by Step

Let's put everything together and solve the open problems!

Initially, each node is itself a cluster of size 1 and no edges are selected.

The algorithm consists of 6 steps. Each step can be performed in constant time.

All 6 steps together build one phase of the algorithm, thus the time complexity of one phase is $O(1)$.

A specific node in each cluster F , e.g. the node with minimum ID, is considered the leader $\ell(F)$ of the cluster F .



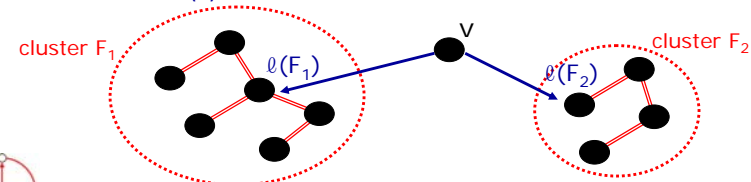
III. Fast MST Algorithm



The Algorithm Step by Step

Step 1

- (a) Each node v computes the minimum-weight edge $e(v, F)$ that connects v to any node of cluster F for all clusters other than the own cluster
- (b) Each node v sends $e(v, F)$ (including the edge weight) to the leader $\ell(F)$ for all clusters other than the own cluster



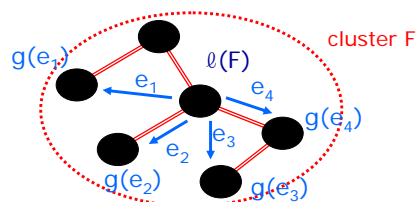
III. Fast MST Algorithm



The Algorithm Step by Step

Step 2

- (a) Each leader v of a cluster F (i.e. $\ell(F) = v$) computes the lightest edge between F and every other cluster
- (b) Each leader v performs procedure Cheap_Out \rightarrow Selects the β lightest outgoing edges and appoints them to its nodes



If edge e is appointed to v , then v is denoted e 's guardian $g(e)$



III. Fast MST Algorithm



The Algorithm Step by Step

Procedure Cheap_Out

Code for the leader of cluster F

Input: Lightest edge $e(F, F')$ for every other cluster F'

1. Sort the input edges in increasing order of weight
2. Define $\beta = \min\{|F|, \text{\# of clusters}\}$
3. Choose the first β edges of the sorted list
4. Appoint the node with the i th largest ID as the guardian of the i th edge, $i = 1, \dots, \beta$
5. Send a message about the edge to the node it is appointed to



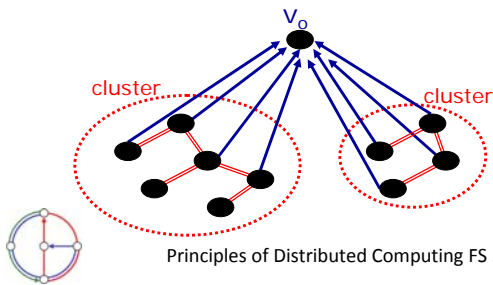
III. Fast MST Algorithm



The Algorithm Step by Step

Step 3

All nodes, that are guardians for a specific edge, send a message to the designated node v_0 , e.g. the node with the smallest ID in the graph



v_0 knows the β lightest outgoing edges of each cluster!

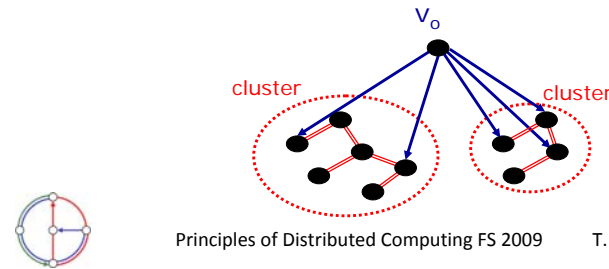
III. Fast MST Algorithm



The Algorithm Step by Step

Step 4

- (a) v_0 locally performs procedure `Const_Frags` \rightarrow Computes the edges to be added
- (b) For all added edges, v_0 sends a message to $g(e)$



III. Fast MST Algorithm

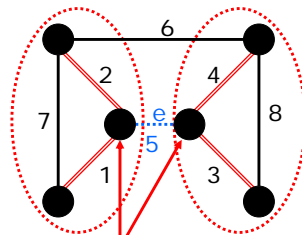


The Algorithm Step by Step

How does `Const_Frags` work?

As we've seen before, a problem occurs when all β outgoing edges of a cluster are used up!

More precisely, a problem occurs only if there is at least one cluster in each of the two super-clusters that are supposed to be merged!!



Used up all edges!

III. Fast MST Algorithm



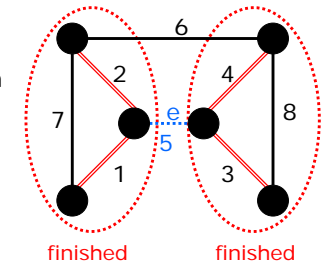
The Algorithm Step by Step

How does `Const_Frags` work?

Finished means **not safe** (see Ch. 10 of the script)

We call a super-cluster containing a cluster that used up all of its β edges **finished**.

If an edge is the lightest outgoing edge of one super-cluster that is **not finished**, then it is still safe to add it, no matter if the other super-cluster is finished, since we are sure that there is no better edge to connect the unfinished super-cluster to other clusters.



III. Fast MST Algorithm



The Algorithm Step by Step

Procedure Const_Frags

Code for the designated node v_0

Input: the β lightest outgoing edges of each cluster

1. Construct the logical graph
2. Sort the input edges in increasing order of weight
3. Go through the list, starting with the lightest edge:
If the edge can be added without creating a cycle then
 add it
else
 drop it



III. Fast MST Algorithm



The Algorithm Step by Step

Procedure Const_Frags

If the two super-clusters are merged, then the new super-cluster is declared finished if

- the edge is the heaviest edge of a cluster in any of the two super-clusters or
- any of the two super-clusters is already finished

If the edge is dropped (\rightarrow both clusters already belong to the same super-cluster) then the super-cluster containing the clusters the edge connects is declared finished if

- the edge is the heaviest edge of any of the two clusters



III. Fast MST Algorithm



The Algorithm Step by Step

Procedure Const_Frags

Note: If a super-cluster is declared finished then it will remain finished

Final step:

All edges between finished super-clusters are **deleted** (before looking at the next lightest edge)

Those are the dangerous edges!



III. Fast MST Algorithm



The Algorithm Step by Step

Step 5

All nodes, that received a message from v_0 , broadcast their edge to all other nodes

Step 6

Each node adds all edges and computes the new clusters

If the number of clusters is greater than 1, then the next phase starts



III. Fast MST Algorithm



The Algorithm Step by Step

The entire algorithm for node v in cluster F

1. Compute the minimum-weight edge $e(v, F')$ that connects v to cluster F' and send it to $\ell(F')$ for all clusters $F' \neq F$
2. if $v = \ell(F)$: Compute lightest edge between F and every other cluster. Perform **Cheap_Out**
3. if $v = g(e)$ for some edge e : Send $\langle e \rangle$ to v_0
4. if $v = v_0$: Perform **Const_Frags**. Send message to $g(e)$ for each added edge e
5. if v received a message from v_0 : Broadcast it
6. Add all received edges and compute the new clusters



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 - Correctness
 - Time and Message Complexity
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IV. Analysis



Correctness

It suffices to show that whenever an edge is added, it is part of the MST \rightarrow We only have to analyze **Const_Frags**!

Proof [Sketch]: We only have to show that we always add the lightest outgoing edge of a super-cluster. Because of Lemma 1, this is always the right choice!



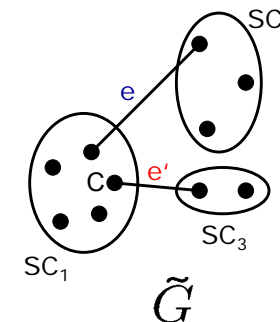
IV. Analysis



Correctness

Assume edge e is used to merge super-cluster SC_1 and SC_2 . W.l.o.g., assume that SC_1 is not finished and that e is one of the β lightest outgoing edges of its cluster. We show that e is the MWOE of SC_1 !

Assume that there is a lighter outgoing edge e' ($w(e') < w(e)$), incident to a cluster C , that connects super-cluster SC_1 to super-cluster SC_3 .



● is a cluster



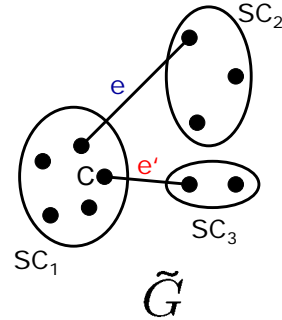
IV. Analysis

Correctness

Case 1: e' is among the β lightest outgoing edges of its cluster C

Since $w(e') < w(e)$, e' must have been considered before e , thus either SC_1 and SC_3 have been merged before or e' was dropped because $SC_1 = SC_3$. Either way, e' cannot be an outgoing edge when the algorithm adds e .

→ Contradiction!



IV. Analysis

Correctness

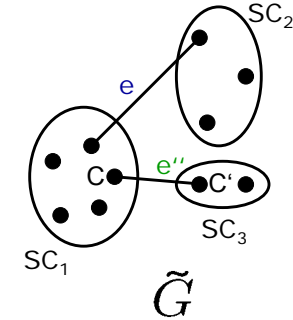
Case 2: e' is not among the β lightest outgoing edges of its cluster C

Case 2.1: There is an edge e'' among the β lightest outgoing edges from C to the same cluster C' :

It follows that $w(e'') < w(e')$. Since $SC_1 \neq SC_3$, e'' has not been considered yet, thus $w(e) < w(e'')$. It follows that

$w(e) < w(e')$.

→ Contradiction!



IV. Analysis

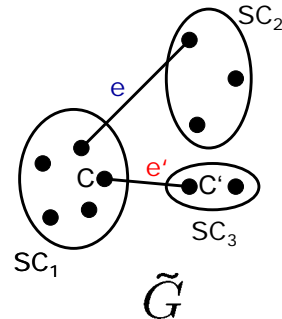
Correctness

Case 2: e' is not among the β lightest outgoing edges of its cluster C

Case 2.2: There is no lighter edge among the β lightest outgoing edges from C to C' :

Thus, all β outgoing edges (of C) have lower weights than e' , also the heaviest of these edges e'' , i.e., $w(e'') < w(e') < w(e)$. Hence, edge e'' must have been inspected already. Since it is the heaviest (last) edge of some cluster, SC_1 must now be finished.

→ Contradiction!



IV. Analysis

Time Complexity

Each phase requires $O(1)$ rounds, but how many phases are required until termination?

Reminder: β_k denotes the **minimum cluster size in phase k** .

Lemma 2: It holds that

$$\beta_{k+1} \geq \beta_k (\beta_k + 1).$$



IV. Analysis

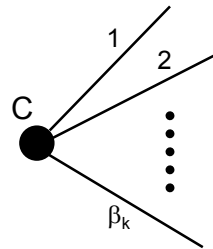


Time Complexity

Proof [Sketch]:

We prove a stronger claim: Whenever a super-cluster is declared finished in phase $k+1$, it contains at least $\beta_k + 1$ clusters.

Each cluster has (at least) β_k outgoing edges in phase $k+1$, since β_k is the minimal cluster size after phase k .



IV. Analysis

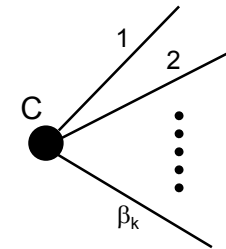


Time Complexity

Case 1: The super-cluster is declared finished after one of its clusters has used up all of its β_k outgoing edges. Let C be a cluster.

Let's call those edges edge 1, 2, ..., β_k leading to the clusters $C_1, C_2, \dots, C_{\beta}$.

If the inspection of an edge does not result in a merge, then the clusters already belong to the same super-cluster! If there is a merge, then they belong to the same super-cluster afterwards.



IV. Analysis

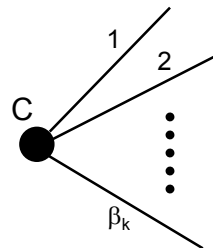


Time Complexity

Thus, at the end, the super-cluster contains at least $C, C_1, C_2, \dots, C_{\beta}$!
→ The super-cluster contains at least $\beta_k + 1$ clusters.

Case 2: The super-cluster is declared finished after merging with an already finished super-cluster.

Using an inductive argument, the finished super-cluster must already contain at least $\beta_k + 1$ clusters, since one of its clusters has used up all of its β_k edges. ■



IV. Analysis



Time Complexity

Theorem 1: The time complexity is $O(\log \log n)$ rounds.

Proof: By Lemma 2, it holds that $\beta_{k+1} \geq \beta_k (\beta_k + 1)$. Furthermore it holds that $\beta_0 := 1$. Hence it follows that

$$\beta_k \geq 2^{2^{k-1}}$$

for every $k \geq 1$. Since $\beta_k \leq n$, it follows that $k \leq \log(\log n) + 1$. Since each phase requires $O(1)$ rounds, the time complexity is $O(\log \log n)$. ■



IV. Analysis



Message Complexity

Theorem 2: The message complexity is
 $O(n^2 \log n)$.

The proof is simple: Count the messages exchanged in steps 1, 3, 4 and 5. We will not do that here.

Adler et al. showed that the minimal number of bits required to solve the MST problem in this model is $\Omega(n^2 \log n)$, thus this algorithm is asymptotically optimal!!!



Overview



- I. Introduction
- II. Previous Results
- III. The New Algorithm
- IV. Analysis
- V. Summary
 - [Results & Conclusions](#)
- VI. Extensive Example



V. Summary



Results

The presented algorithm solves the MST problem in the given model in $O(\log \log n)$ rounds.

The algorithm sends $O(n^2 \log n)$ bits in total, which is asymptotically optimal.



V. Summary



Conclusions

„An obvious question we leave open is whether the algorithm can be improved, or is there an inherent lower bound of $\Omega(\log \log n)$ on the number of communication rounds required to construct an MST in this model.“

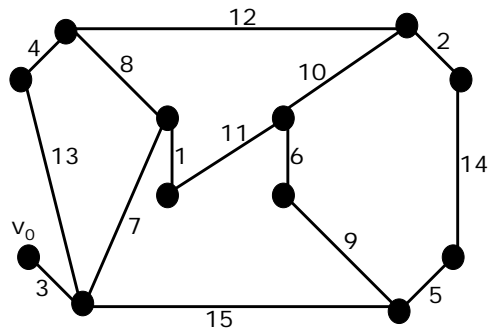
- Is there a faster algorithm?
- Is $\Omega(\log \log n)$ a lower bound?
- Is there an algorithm with time complexity $O(\log \log n)$ for graphs of diameter 2?



VI. Extensive Example



The graph



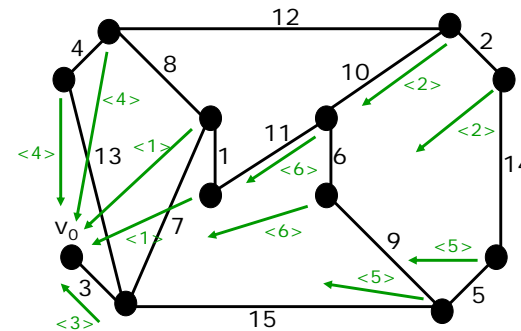
All other edges
are heavier!!!



VI. Extensive Example



Phase 1



1. Not necessary
2. Not necessary
3. Send MWOE to v_0
4. Const_Frags!

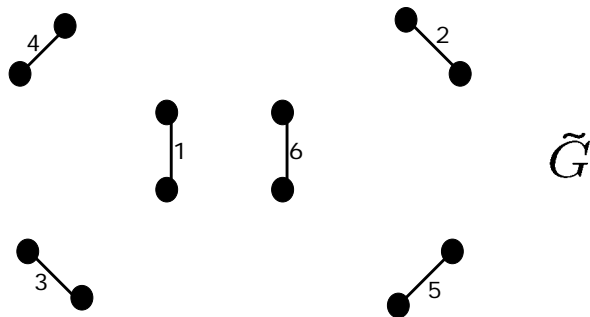


VI. Extensive Example



Phase 1 – Const_Frags

1. Construct logical graph

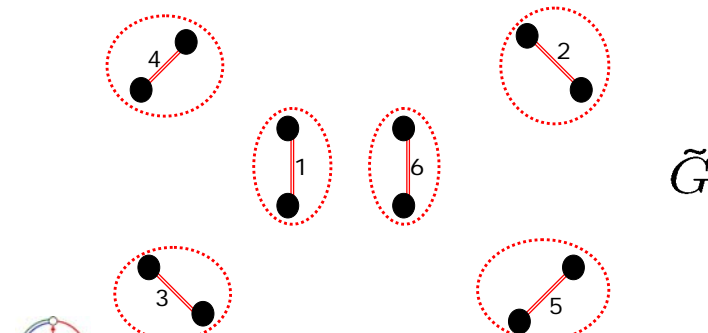


VI. Extensive Example



Phase 1 – Const_Frags

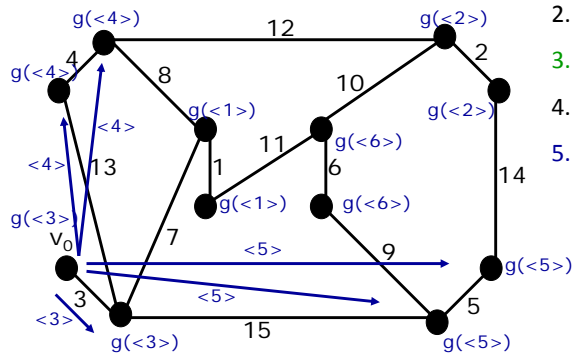
2. Add edges



VI. Extensive Example



Phase 1



1. Not necessary
2. Not necessary
3. Send MWOE to v_0
4. Const_Frags!
5. Send e to $g(e)$

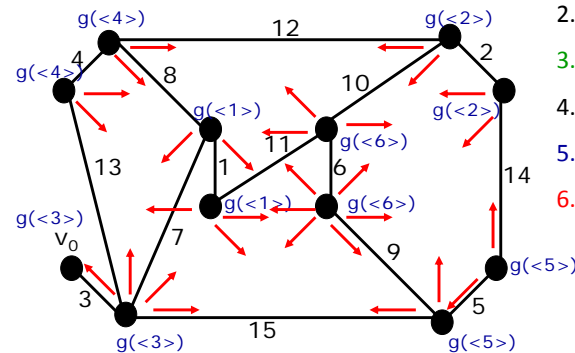
Only some messages are displayed!



VI. Extensive Example



Phase 1



1. Not necessary
2. Not necessary
3. Send MWOE to v_0
4. Const_Frags!
5. Send e to $g(e)$
6. Broadcast e and update the clusters

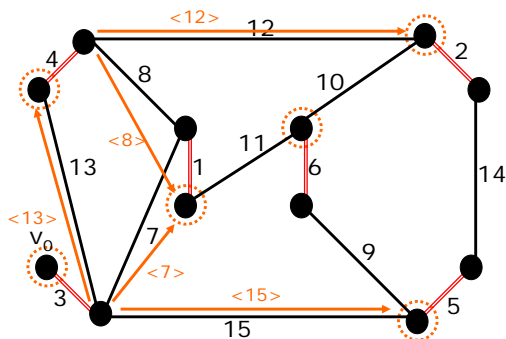
Only some messages are displayed!



VI. Extensive Example



Phase 2



1. Compute $e(v, F')$ and send it to $\ell(F')$

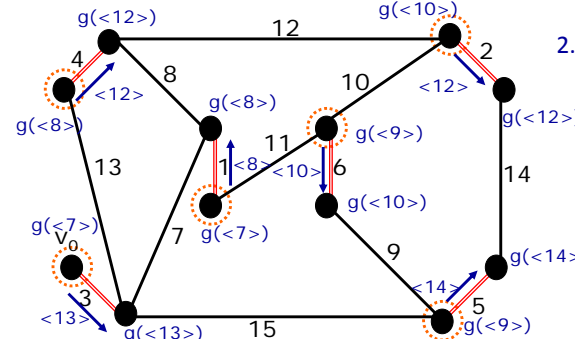
Only some messages are displayed!



VI. Extensive Example



Phase 2



1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $\beta = 2$ lightest outgoing edges and appoint guardians

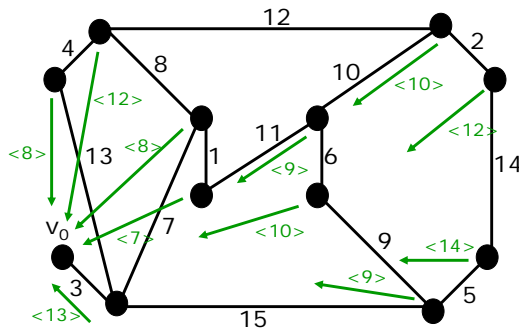
Only some messages are displayed!



VI. Extensive Example



Phase 2



1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $b = 2$ lightest outgoing edges and appoint guardians
3. Send appointed edge to v_0
4. Const_Frags!

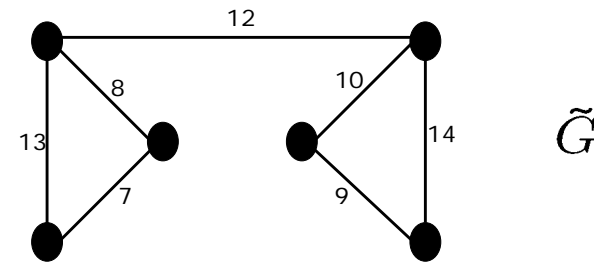


VI. Extensive Example



Phase 2 – Const_Frags

1. Construct logical graph

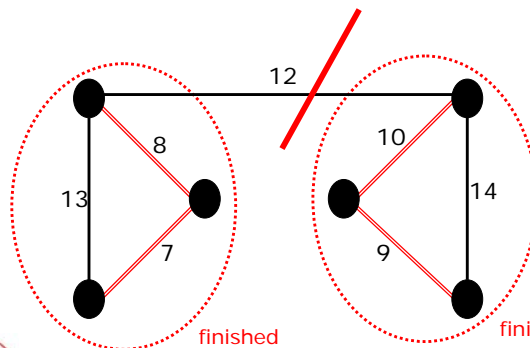


VI. Extensive Example



Phase 2 – Const_Frags

1. Add edges



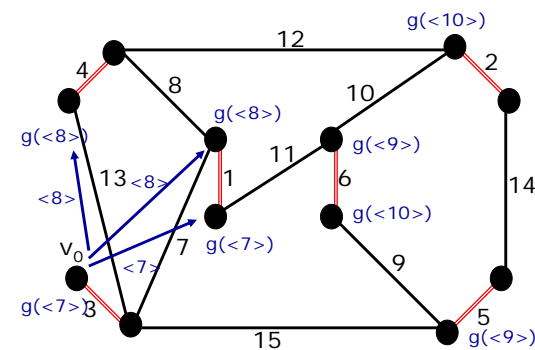
Edges between finished super-clusters must not be added!



VI. Extensive Example



Phase 2



1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $b = 2$ lightest outgoing edges and appoint guardians
3. Send appointed edge to v_0
4. Const_Frags!
5. Send e to $g(e)$

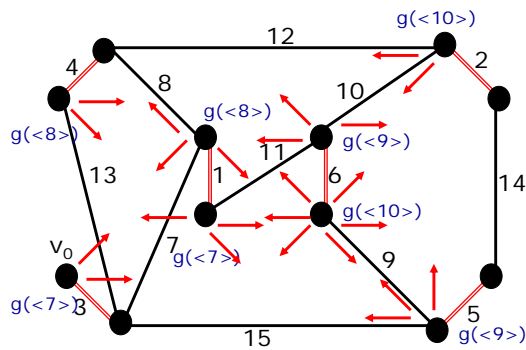
Only some messages are displayed!



VI. Extensive Example



Phase 2



Only some messages are displayed!

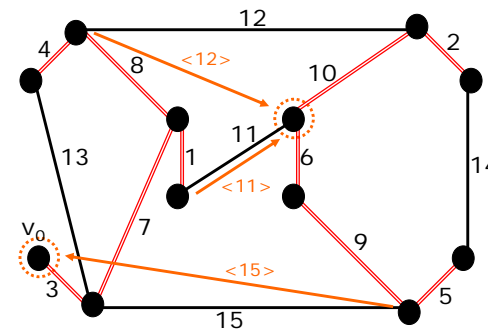
1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $b = 2$ lightest outgoing edges and appoint guardians
3. Send appointed edge to v_0
4. Const_Frags!
5. Send e to $g(e)$
6. Broadcast e and update the clusters



VI. Extensive Example



Phase 3



Only some messages are displayed!

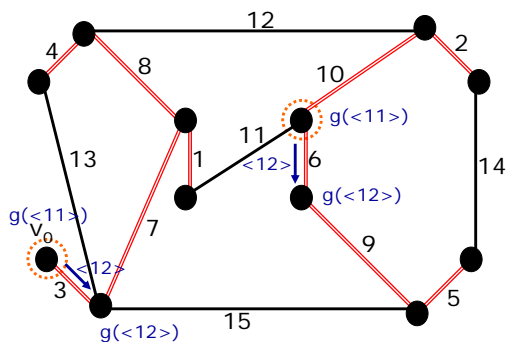
1. Compute $e(v, F')$ and send it to $\ell(F')$



VI. Extensive Example



Phase 3



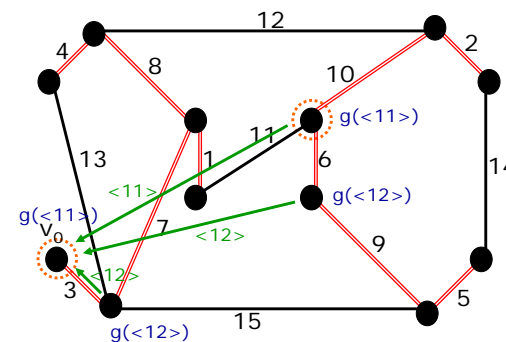
1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $\beta = 2$ lightest outgoing edges and appoint guardians



VI. Extensive Example



Phase 3



1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $\beta = 2$ lightest outgoing edges and appoint guardians
3. Send appointed edge to v_0
4. Const_Frags!

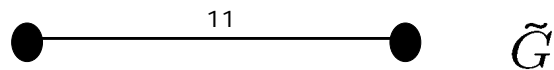


VI. Extensive Example



Phase 3 – Const_Frags

1. Construct logical graph



VI. Extensive Example



Phase 3 – Const_Frags

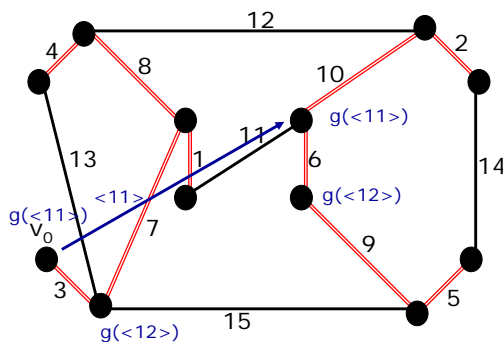
1. Add edges



VI. Extensive Example



Phase 3



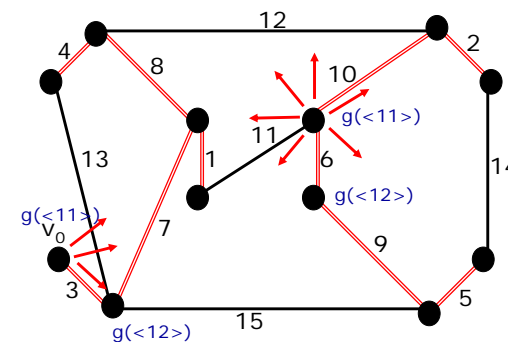
1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $\beta = 2$ lightest outgoing edges and appoint guardians
3. Send appointed edge to v_0
4. Const_Frags!
5. Send e to $g(e)$



VI. Extensive Example



Phase 3



1. Compute $e(v, F')$ and send it to $\ell(F')$
2. Select $b = 2$ lightest outgoing edges and appoint guardians
3. Send appointed edge to v_0
4. Const_Frags!
5. Send e to $g(e)$
6. Broadcast e and update the clusters

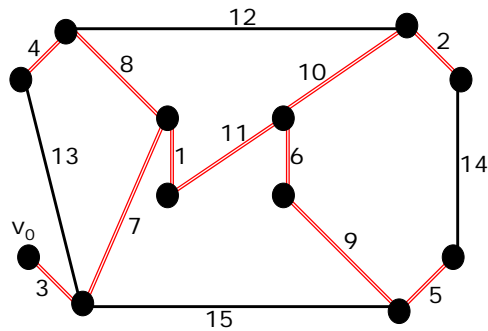
Only some messages are displayed!



VI. Extensive Example



After phase 3



Done! 😊



References



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